#### Pregel: A System for Large-Scale Graph Processing

Written by G. Malewicz et al. at SIGMOD 2010 Presented by Chris Bunch Tuesday, October 12, 2010





## Graphs are hard

Poor locality of memory access

- Very little work per vertex
- Changing degree of parallelism
- Running over many machines makes the problem worse





## State of the Art Today

Write your own infrastructure
Substantial engineering effort
Use MapReduce

 Inefficient - must store graph state in each stage, too much communication between stages





## State of the Art Today

Use a single-computer graph library
Not scalable :
Use existing parallel graph systems
No fault tolerance :





# Bulk Synchronous Parallel

- Series of iterations (supersteps)
- Each vertex invokes a function in parallel
- Can read messages sent in previous superstep
- Can send messages, to be read at the next superstep
- Can modify state of outgoing edges





## Compute Model

- You give Pregel a directed graph
- It runs your computation at each vertex
- Do this until every vertex votes to halt
- Pregel gives you a directed graph back



#### Primitives

- Vertices first class
- Edges not
- Both can be dynamically created and destroyed





## Vertex State Machine







- Your code subclasses Vertex, writes a Compute method
- Can get/set vertex value
- Can get/set outgoing edges values
- Can send/receive messages





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- Message passing:
- No guaranteed message delivery order
- Messages are delivered exactly once
- Can send messages to any node
- If dest doesn't exist, user's function is called





- Combiners (off by default):
- User specifies a way to reduce many messages into one value (ala Reduce in MR)
- Must be commutative and associative
- Exceedingly useful in certain contexts (e.g., 4x speedup on shortest-path compution)





- Aggregators:
- User specifies a function
- Each vertex sends it a value
- Each vertex receives aggregate(vals)
- Can be used for statistics or coordination





- Topology mutations:
- Vertices can create / destroy vertices at will
- Resolving conflicting requests:
  - Partial ordering: E Remove, V Remove, V Add, E Add
  - User-defined handlers: You fix the conflicts on your own





• Input and output:

- Text file
- Vertices in a relational DB
- Rows in BigTable
- Custom subclass Reader/Writer classes





#### Implementation

Executable is copied to many machines
One machine becomes the Master
Coordinates activities
Other machines become Workers
Performs computation





#### Implementation

- Master partitions the graph
- Master partitions the input
  - If a Worker receives input that is not for their vertices, they pass it along
- Supersteps begin
- Master can tell Workers to save graphs





#### Fault Tolerance

- At each superstep S:
  - Workers checkpoint V, E, and Messages
  - Master checkpoints Aggregators
- If a node fails, everyone starts over at S
- Confined recovery is under development
- what happens if the Master fails?



## The Worker

- Keeps graph in memory
- Message queues for supersteps S and S+I
- Remote messages are buffered
- Combiner is used when messages are sent or received (save network and disk)





### The Master

- Master keeps track of which Workers own each partition
  - Not who owns each Vertex
- Coordinates all operations (via barriers)
- Maintains statistics and runs a HTTP server for users to view info on





## Aggregators

Worker passes values to its aggregator

- Aggregator uses tree structure to reduce vals w/ other aggregators
  - Better parallelism than chain pipelining
- Final value is sent to Master





# PageRank in Pregel

```
class PageRankVertex
    : public Vertex<double, void, double> {
public:
  virtual void Compute(MessageIterator* msgs) {
    if (superstep() >= 1) {
      double sum = 0;
      for (; !msgs->Done(); msgs->Next())
        sum += msgs->Value();
      *MutableValue() =
          0.15 / NumVertices() + 0.85 * sum;
    }
    if (superstep() < 30) {</pre>
      const int64 n = GetOutEdgeIterator().size();
      SendMessageToAllNeighbors(GetValue() / n);
    } else {
      VoteToHalt();
    3
  ጉ
};
```





## Shortest Path in Pregel

```
class ShortestPathVertex
    : public Vertex<int, int, int> {
  void Compute(MessageIterator* msgs) {
    int mindist = IsSource(vertex_id()) ? 0 : INF;
    for (; !msgs->Done(); msgs->Next())
      mindist = min(mindist, msgs->Value());
    if (mindist < GetValue()) {</pre>
      *MutableValue() = mindist;
      OutEdgeIterator iter = GetOutEdgeIterator();
      for (; !iter.Done(); iter.Next())
        SendMessageTo(iter.Target(),
                      mindist + iter.GetValue());
    VoteToHalt();
};
```





#### Evaluation

- 300 multicore commodity PCs used
- Only running time is counted
  - Checkpointing disabled
- Measures scalability of Worker tasks
- Measures scalability w.r.t. # of Vertices
  - in binary trees and log-normal trees







Figure 7: SSSP—1 billion vertex binary tree: varying number of worker tasks scheduled on 300 multicore machines







Figure 8: SSSP—binary trees: varying graph sizes on 800 worker tasks scheduled on 300 multicore machines





Figure 9: SSSP—log-normal random graphs, mean out-degree 127.1 (thus over 127 billion edges in the largest case): varying graph sizes on 800 worker tasks scheduled on 300 multicore machines





## Current / Future Work

- Graph must fit in RAM working on spilling over to / from disk
- Assigning vertices to machines to optimize traffic is an open problem
  - Want to investigate dynamic repartitioning





#### Conclusions

- Pregel is production-ready and in use
- Usable after a short learning curve
  - Vertex centric is not always easy to do
- Pregel works best on sparse graphs w / communication over edges
- Can't change the API too many people using it!





## Related Work



- Hama from the Apache Hadoop team
- BSP model but not vertex centric ala Pregel

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Appears not to be ready for real use:
 Hama Wiki Login
 GettingStarted

FrontPage RecentChanges FindPage HelpContents GettingStarted

Immutable Page Info Attachments More Actions:

#### NOTE: Hama is not ready yet!!





## Related Work

Phoebus, released last week on github
Runs on Mac OS X
Cons (as of this writing):
Doesn't work on Linux
Must write code in Erlang (since Phoebus is written in it)





## Thanks!

To my advisor, Chandra Krintz
To Google for this paper
To all of you for coming!



